Section 7.1

X: the consumption set

Notation for preferences: $x \succeq y$ or $\mathbf{x} \succeq \mathbf{y}$. Also, $\mathbf{x} \succ \mathbf{y}$, $\mathbf{x} \preceq \mathbf{y}$, $\mathbf{x} \prec \mathbf{y}$, or $\mathbf{x} \sim \mathbf{y}$. (No prices nor income; not a market environment; psychology only.)

A consumer is "rational" if preferences are:

- "complete": assuming x ≠ y, either x ≻ y, or y ≻ x, or x ~ y. Implications: no learning. Difficult example: choose which of your children to give up.
- ["reflexive": Varian says this is needed but it's not.]
- "transitive": if $x \succeq y$ and $y \succeq z$ then $x \succeq z$. (Sometimes violated.)

"Continuity" of preferences: suppose an infinite sequence \mathbf{x}_i is convergent and call its limit \mathbf{x}^* . If $\mathbf{x}_i \succeq \mathbf{y}$ for all *i*, then "continuity of preferences" requires $\mathbf{x}^* \succeq \mathbf{y}$.

A theorem (MCWG p. 47): if a consumer's preferences are "rational" and "continuous" then those preferences can be represented by a continuous function mapping X into \mathbb{R}^1 . In other words, there exists at least one function $u(\mathbf{x}) : X \to \mathbb{R}^1$ which satisfies

$$\mathbf{x} \succ \mathbf{y} \Leftrightarrow u(\mathbf{x}) > u(\mathbf{y}) \quad \forall \mathbf{x}, \mathbf{y} \in X.$$

We call this function a "utility function." Cf. Varian p. 97.

First problem with these assumptions: perhaps preferences, instead of being fixed, depend on the question asked. See MCWG p. 7:

Consider the following example, paraphrased from Kahneman and Tversky (1984):

Imagine that you are about to purchase a stereo for 125 dollars and a calculator for 15 dollars. The salesman tells you that the calculator is on sale for 5 dollars less at the other branch of the store, located 20 minutes away. The stereo is the same price there. Would you make the trip to the other store?

It turns out that the fraction of respondents saying that they would travel to the other store for the 5 dollar discount is much higher than the fraction who say they would travel when the question is changed so that the 5 dollar saving is on the stereo. This is so even though the ultimate saving obtained by incurring the inconvenience of travel is the same in both cases. Second problem with these assumptions: people may not know what makes them happy.

- 1. Daniel Gilbert, Harvard Psychology Dept.: discussed in Sept. 7, 2003 New York Times, "The Futile Pursuit of Happiness." Happiness set points.
- 2. Baba Shiv https://whywereason.wordpress.com/tag/baba-shiv/. Cognitive processing is hard; the brain is not monolithic.

Third problem with these assumptions: lexicographic preferences (MCWG p. 46) violate the continuity assumption. Suppose a consumer always prefers bundles having more chocolate to those having less chocolate regardless of what else is in the bundles. For example, if chocolate is the second element in the consumption vector, this consumer would have:

$$(1, 1 + \frac{1}{i}) \succ (2, 1)$$
 for all $i > 0$.

These preferences are lexicographic, and while the above ranking makes sense for all $i < \infty$, it makes no sense in the limit as $i \to \infty$ —in other words, the limiting bundle, $\mathbf{x}^* = (1, 1)$, does not satisfy $(1, 1) = \mathbf{x}^* \succ (2, 1)$ so these preferences violate continuity. (Mention Nicholas Georgescu-Roegen.)

Common assumptions on preferences:

- local nonsatiation
- weak monotonicity ("more is never worse")
- strong monotonicity ("more is strictly better")

Historical example of violation of monotonicity: the "potlatch" of the Native Americans of the Pacific Northwest.

Claim: strong monotonicity implies local nonsatiation.



Claim: local nonsatiation does not imply strong monotonicity. [In the counterexample, the straight lines are "indifference curves," which are defined to be the contour lines ("level sets") of the utility function, and also, $U_0 < U_1 < U_2 < U_3$.]



"Convex" or "strictly convex" preferences. Do not confuse this "convexity" with:

- convex combinations (of, especially, vectors)
- convex sets
- convex functions.

"Convex preferences" are a type of convex binary relation: roughly (for the exact definitions see p. 96 of Varian),

 $\mathbf{x} \succeq \mathbf{z}$ and $\mathbf{y} \succeq \mathbf{z} \implies \forall t \in (0, 1), t \mathbf{x} + (1-t) \mathbf{y} \begin{cases} \succeq \mathbf{z} & \text{convex preferences} \\ \succ \mathbf{z} & \text{strictly convex preferences.} \end{cases}$

A graph which shows preferences that are convex but not strictly convex (the straight line is an indifference curve):



Claim: if preferences are strictly convex and if $\mathbf{x} \sim \mathbf{y}$, then

$$t \mathbf{x} + (1 - t) \mathbf{y} \succ \mathbf{x} \sim \mathbf{y}$$
.

Recall that from Varian's (rigorously correct) definition of strictly convex preferences, in the special case of the diagram below, " $\mathbf{x} \neq \mathbf{y}, \mathbf{x} \succeq \mathbf{z}, \mathbf{y} \succeq \mathbf{z}$ " implies $t \mathbf{x} + (1 - t) \mathbf{y} \succ \mathbf{z} = \mathbf{y} \sim \mathbf{x}$. (Hence there is a relation between "preferences are a strictly convex binary relation" and "indifference curves are a strictly convex function of, in this graph, the single variable 'apples'.")



EXistence of U: MWGp.47 (retional: complete of transitive (Virian) P.2 MWG. P. 6 too

Marginel Rate of Substitution: du(x) = 0 or, butter: u(x) = constant2u(x)

$$\frac{\partial u(x)}{\partial x_{i}} dx_{i} + \frac{\partial u(x)}{\partial x_{j}} dx_{j} = 0$$

$$x_{i}$$

$$\frac{dx_{j}}{dx_{i}} = -\frac{\partial u/\partial x_{i}}{\partial u/\partial x_{j}}$$

$$\int u = 0 \quad \text{interacting}$$

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$$\frac{dv}{du} \frac{\partial u}{\partial x_{i}} x_{i} + \frac{dv}{du} \frac{\partial u}{\partial x_{j}} dx_{j} = 0$$

$$\frac{dv}{du} \frac{\partial u}{\partial x_{i}} x_{i} + \frac{dv}{du} \frac{\partial u}{\partial x_{j}} dx_{j} = 0$$

Prop. If u(x) represents the preferences of a consumer Not "positive and if v(u) is a strictly moreasing huction of u monotoriz transformation If then v (u(x)) also represents the preferances of that consumer. $v(u(x_1)) > v(u(x_2)) \Rightarrow u(x_1) > u(x_2) \Rightarrow \chi_1 \neq \chi_2.$ Proof. we have the time we have a conce benved from H& D? . Q - Concavity vs. Concovity. v(u)=4 destroys concauty but not monotomeity. v(u)=4 destroys concauty but not Altruism, jealousy. Value of jorny (potlatch) with whether with a way shot to hoppiness.

$$B \triangleq \{ x : x \in X \text{ and } p \colon x \leq m \}$$

max
$$u(x)$$
 s.t. $x \in B$.
Continuous Closed & budd. (p.505)

indirect utility hustion v(p,m) = max u(x) s.t. p.x = mdemand direction x(p,m) - homo geneous of degree <math>On(p,m)

$$\begin{aligned} \lambda &= u(\chi) - \lambda (p \cdot \chi - m) \\ \frac{\partial \chi}{\partial \chi_{c}} &= 0 \implies \frac{\partial u}{\partial \chi_{c}} - \lambda p_{i} = 0 \end{aligned}$$



 $f(x + \Delta x) \approx f(x) + \nabla f(x) \cdot \Delta x + \frac{1}{2} \Delta x \quad \nabla^2 f \Delta x$ $f(x + \Delta x) - f(x) = \nabla f \cdot \Delta x + \frac{1}{2} \Delta x^{T} \nabla^2 f \Delta x \leq 0?$ $\int_{U}^{U} \int_{O}^{U} \int_{O}^{U} f \int_$

(7,2)

Indirect utility function & (R. m) to prove, look at the 1) nonnereasing in p : nondecreasing in m attordable set 2) homogeneous of depres O a (P, m) 3) grassiconvex in P $\{p: \nu(p, m) \leq k\}$ utility of k (proof united) P. 4) continuous (p.506 Them, (1).) Expenditure function e(p, u) = mon p. x s.t. u(x) >, U 1) nondecreasing & R: 2 Pri= hi 70 (Shephand's Lemma) 20 most page fint! -- e(tp,+(1-+)p) 2) homoj. degree 1 in P min >p·x = ... "" "notation 3) concome in p (pour destrict) expanditure dumb consumer P. X, +... =min(tp,+(1-1)p)x smart consumer tp. +(1++)pz min (1-t) pz x Pi Don't Confusp 4) continuous in of for proof omitted) and 1 5) Hicksian demand function h (p, u) is the answer to the exp. min. problem. Max(ftg) vs. maxft max g Shephard's Lemma: $h:(p, u) = \frac{\partial e(p, u)}{\partial p_i}$, Proof -> Hom? -7 Marshalloan is × (p, m). MC WIG : "Walresign"

$$Envelope Theorem (p.50f)$$

$$M(a) = max g(x_1, x_2, a)$$

$$s.t.$$

$$h(x_1, x_2, a) = 0.$$

$$d M(a) = \frac{\partial x}{\partial a}$$

7.3 p.2

Prop. Utility maximization => expenditive minimization.

Assume local unsatiation

Let
$$\chi^{*u}$$
 solve max $u(\chi)$ s.t. $p \cdot \chi \le m$.
Let $\hat{u} = u(\chi^{*u})$.
Then χ^{*u} solves min $p \cdot \chi$ s.t. $u(\chi) \neq \hat{u}$.
Proof. Suppose not. Let χ^{*e} solve min $p \cdot \chi$ s.t. $u(\chi) \neq \hat{u}$. Note that χ^{*u} is advantille
Then $p \cdot \chi^{*e} and $u(\chi^{*e}) \neq \hat{u} = u(\chi^{*u})$.
Let χ^{*u} solves then χ^{*e} would solve the utility -meximize ation
we'd be done be cause then χ^{*e} would solve the utility -meximize ation
problem, what can treat to solve assumption that χ^{*u} is blue that problem.
How about if $u(\chi^{*e}) = u(\chi^{*u})$?
Then by local humschiation $\exists \chi''$ close enough to χ^{*e} that it still
satisfies $p \cdot \chi'' , but χ'' improves on χ^{*e} :
 $u(\chi'') > u(\chi^{*e}) = u(\chi^{*u})$.
Then χ'' , not χ^{*u} , solves the utility -meximized time problem, a
low factoring.$$

7,4 p.2

Prop. Expendifive mainimization
$$\Rightarrow$$
 utility maximization.
Assume containing of u.
Let χ^{*e} solve min $p \cdot \chi$ st. $u(\chi) > \hat{u}$.
Let $m = p \cdot \chi^{*e}$.
Then χ^{*e} solves max $u(\chi)$ s.t. $p \cdot \chi \le m$.
Proof. Supposend, and instead let χ^{*u} solve max $u(\chi)$ s.t. $p \cdot \chi \le m$. Note that
then $u(\chi^{*u}) > u(\chi^{*e})$ and $p \cdot \chi^{*u} \le m = p \cdot \chi^{*e}$.
By the definition of χ^{*e} , $u(\chi^{*e}) > \hat{u}$, so
 $u(\chi^{*u}) > u(\chi^{*e}) > \hat{u}$.
[If $p \cdot \chi^{*u}
we'd be done because then χ^{*u} would solve the expendence multivatorial
problem, contradicting our assumption that χ^{*e} does so.]
Let $\chi \in (o,1)$ (thick of it being very close to 1). Then
 $p \cdot (\chi^{*u}) > u(\chi^{*u}) > u(\chi^{*u}) = p \cdot \chi^{*u} \le p \cdot \chi^{*e}$.
and by continuity of $u(\cdot)$, since $u(\chi^{*u}) > u(\chi^{*e})$, for λ close to 1.
Then $\lambda \chi^{*u}$ solves the expenditive multivatorial multivatorial
 $u(\chi^{*u}) > u(\chi^{*u}) = \mu(\chi^{*u}) > u(\chi^{*e})$.$

7,4 p.3

Roy's Identity:
$$X_i(p,m) = -\frac{\partial v(p,m)/\partial P_i}{\partial v(p,m)/\partial m}$$
 for $i = 1, 2, ..., k$.

$$\frac{Proof}{Envelope thm.}$$

$$M(a) = max g(x_1, x_2, a)$$

$$s.t.$$

$$t(x_1, x_2, a) = 0.$$

$$\begin{aligned} \mathcal{L} &= g - \lambda h \\ \frac{\partial M(a)}{\partial a_i} &= \frac{\partial \mathcal{L}^*}{\partial a_i} \end{aligned}$$

Indirect Whileby Function.

$$v(p,m) =$$

 $\max_{x} u(x)$
 x
s.t.

$$p \cdot \chi - m = O$$

$$\mathcal{L} = u(\mathbf{x}) - \lambda \left[\mathbf{p} \cdot \mathbf{x} - \mathbf{m} \right]$$

$$\begin{cases} \frac{\partial v}{\partial m} = \frac{\partial}{\partial m} (\mathcal{L}^*) = \lambda \quad (sensitivity analysis) \\ \frac{\partial v}{\partial p_i} = \frac{\partial}{\partial p_i} (\mathcal{L}^*) = -\lambda \times i \\ substituting for \lambda, one has \end{cases}$$

$$\frac{\partial v}{\partial p_i} = - \frac{\partial v}{\partial m} X_i$$
.